

A totally unimodular view of structured sparsity

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Joint work with: Volkan Cevher



Supervised learning and inverse problems

Running example:

The diagram illustrates the equation $b = Ax + w$. On the left is a vertical column vector b with 10 colored squares. In the middle is a matrix A with 10 rows and 10 columns of colored squares, labeled $n \times p$ below it. To the right of A is a vertical column vector x with 10 colored squares. To the right of x is a plus sign $+$. To the right of the plus sign is another vertical column vector w with 10 colored squares. An equals sign $=$ is placed between b and A .

Applications: Machine learning, signal processing, theoretical computer science...

Supervised learning and inverse problems

Running example:

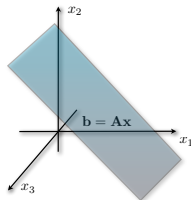
$$\mathbf{b} = \mathbf{A} \mathbf{x}^h + \mathbf{w}$$

$n \times p$

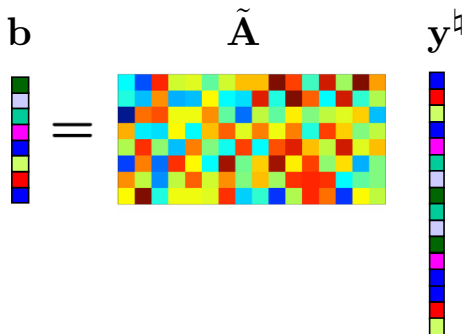
A difficult estimation challenge when $n < p$:

Nullspace (null) of \mathbf{A} : $\mathbf{x}^h + \delta \rightarrow \mathbf{b}, \quad \forall \delta \in \text{null}(\mathbf{A})$

- ▶ **Needle in a haystack:** *We need additional information on \mathbf{x}^h !*

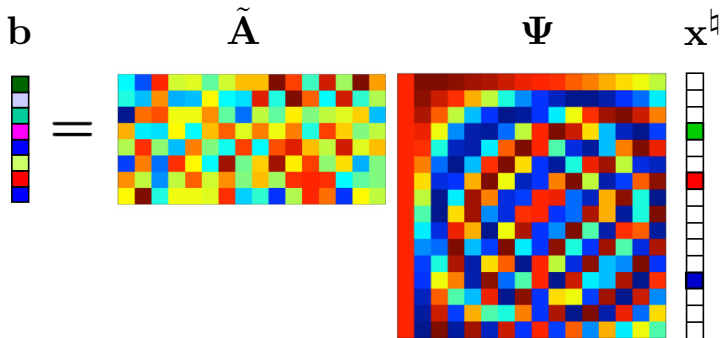


Sparsity to the rescue!

$$\mathbf{b} = \tilde{\mathbf{A}} \mathbf{y}^{\natural}$$


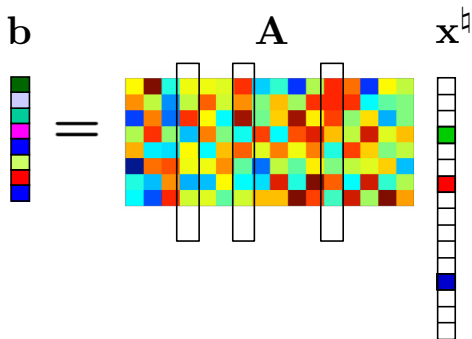
- ▶ $\mathbf{b} \in \mathbb{R}^n$, $\tilde{\mathbf{A}} \in \mathbb{R}^{n \times p}$, and $n < p$

Sparsity to the rescue!



- ▶ $\mathbf{b} \in \mathbb{R}^n$, $\tilde{\mathbf{A}} \in \mathbb{R}^{n \times p}$, and $n < p$
- ▶ $\Psi \in \mathbb{R}^{p \times p}$, $\mathbf{x}^{\natural} \in \Sigma_s$, and $s < n < p$

Sparsity to the rescue!



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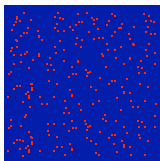
$n \times 1$ $n \times s$ $s \times 1$

- ▶ $\mathbf{b} \in \mathbb{R}^n$, $\mathbf{A} \in \mathbb{R}^{n \times p}$, $\mathbf{x}^{\natural} \in \Sigma_s$, and $s < n < p$

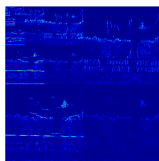
Impact: Support restricted columns of \mathbf{A} leads to an *overcomplete* system.

Beyond sparsity towards model-based or *structured* sparsity

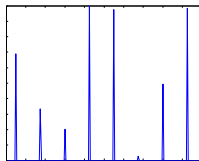
- ▶ The following signals can look the **same** from a **sparsity** perspective!



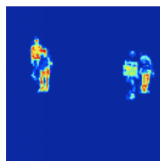
Sparse image



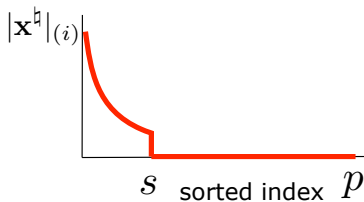
Wavelet coefficients
of a natural image



Spike train

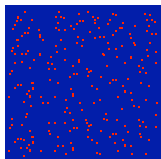


Background subtracted
image

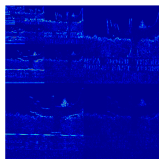


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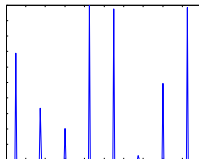
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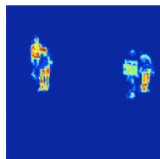
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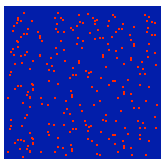


Spike train

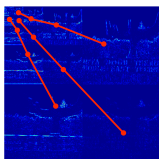


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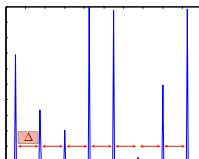
- ▶ In reality, these signals have additional **structures** beyond the simple sparsity



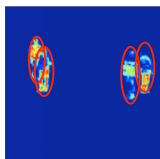
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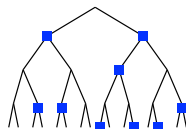
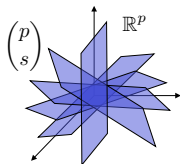
Spike train



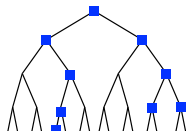
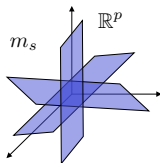
Background subtracted
image

Beyond sparsity towards model-based or *structured* sparsity

Sparsity model: Union of **all** s -dimensional canonical subspaces.



Structured sparsity model: A **particular** union of m_s s -dimensional canonical subspaces.



Three upshots of structured sparsity:

1. Reduced sample complexity
2. Better noise robustness
3. better interpretability

A simple template for linear inverse problems

Find the “*sparsest*” \mathbf{x} subject to *structure* and *data*.

- ▶ **Sparsity**

We can generalize this desideratum to other notions of simplicity

- ▶ **Structure**

We only allow certain sparsity patterns

- ▶ **Data fidelity**

We have many choices of convex constraints & losses to represent data; e.g.,

$$\|\mathbf{b} - \mathbf{A}\mathbf{x}\|_2 \leq \kappa$$

Simple sparsity

A combinatorial approach for estimating \mathbf{x}^{\natural} from $\mathbf{b} = \mathbf{A}\mathbf{x}^{\natural} + \mathbf{w}$

$$\mathbf{x}^{\star} \in \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \|\mathbf{x}\|_0 : \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_2 \leq \|\mathbf{w}\|_2 \right\} \quad (\mathcal{P}_0)$$

where $\|\mathbf{x}\|_0 := \mathbf{1}^T \mathbf{s}$, $\mathbf{s} = \mathbf{1}_{\text{supp}(\mathbf{x})}$, $\text{supp}(\mathbf{x}) = \{i | x_i \neq 0\}$

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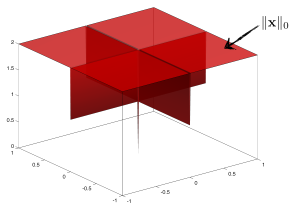
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\mathcal{P}_0 has the following characteristics:

- ▶ sample complexity: $\mathcal{O}(s)$
- ▶ computational effort: NP-Hard
- ▶ stability: No

$\|\mathbf{x}\|_0$ over the unit ℓ_{∞} -ball



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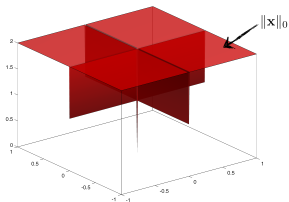
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Convex relaxation:

Convex envelope is the **largest** convex lower bound.

$\|\mathbf{x}\|_0$ over the unit ℓ_{∞} -ball



A technicality: Restrict $\mathbf{x}^{\natural} \in [-1, 1]^p$.

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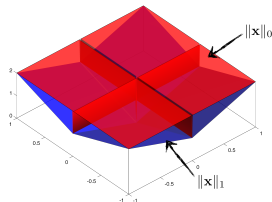
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The role of convexity: Tractable & stable recovery

A combinatorial approach for estimating \mathbf{x}^{\natural} from $\mathbf{b} = \mathbf{A}\mathbf{x}^{\natural} + \mathbf{w}$

$$\mathbf{x}^{\star} \in \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \|\mathbf{x}\|_1 : \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_2 \leq \|\mathbf{w}\|_2, \|\mathbf{x}\|_{\infty} \leq 1 \right\} \quad (BP)$$

where $\|\mathbf{x}\|_1 := \mathbf{1}^T |\mathbf{x}|$

The role of convexity: Tractable & stable recovery

A combinatorial approach for estimating \mathbf{x}^\natural from $\mathbf{b} = \mathbf{A}\mathbf{x}^\natural + \mathbf{w}$

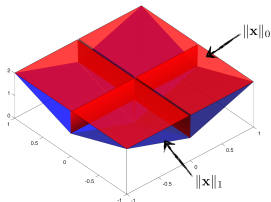
$$\mathbf{x}^* \in \arg \min_{\mathbf{x} \in \mathbb{R}^p} \left\{ \|\mathbf{x}\|_1 : \|\mathbf{b} - \mathbf{A}\mathbf{x}\|_2 \leq \|\mathbf{w}\|_2, \|\mathbf{x}\|_\infty \leq 1 \right\} \quad (BP)$$

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$\|\mathbf{x}\|_1$ is the **convex envelope** of $\|\mathbf{x}\|_0$

BP has the following characteristics [13]:

- ▶ sample complexity: $\mathcal{O}(s \log(\frac{p}{s}))$
- ▶ computational effort: **Tractable**;
 $\mathcal{O}(n^2 p^{1.5} \log(\frac{1}{\epsilon}))$ via IPM (for $w = 0$)
- ▶ stability: Robust to noise



A technicality: Restrict $\mathbf{x}^\natural \in [-1, 1]^p$.

Convex relaxations in general ?

We encode the structure over the support by $g(\mathbf{x}) = F(\text{supp}(\mathbf{x}))$

- ▶ $\text{supp}(\mathbf{x}) = \{i | x_i \neq 0\}$
- ▶ $F(\mathbf{s}) : \{0, 1\}^p \rightarrow \mathbb{R} \cup \{+\infty\}$

How to compute the convex relaxation of g in general ?

1. Case by case heuristics
2. **Biconjugation** (\equiv convex envelope): Fenchel conjugate of Fenchel conjugate.

Recall **Fenchel conjugate**: $g^*(\mathbf{y}) := \sup_{\mathbf{x}:\text{dom}(g)} \mathbf{x}^T \mathbf{y} - g(\mathbf{x})$

Proposition (Hardness of conjugation)

The Fenchel conjugate of g results in the following combinatorial problem

$$g^*(\mathbf{y}) = \sup_{\mathbf{s} \in \{0,1\}^p} |\mathbf{y}|^T \mathbf{s} - F(\mathbf{s}).$$

*which is **NP-Hard** in general.*

Tractable convex relaxation

Prior work:

1. Monotone submodular penalties [1]
 - ▶ Tractable biconjugation via Lovász extension
 - ▶ Limited to certain structures
2. ℓ_q -regularized combinatorial functions [11] ($\mu F(\text{supp}(\mathbf{x})) + \nu \|\mathbf{x}\|_q$)
 - ▶ Tractable biconjugation even for some non-submodular functions
 - ▶ Not always tractable
 - ▶ May loose structure

Our work: New framework for tractable convex relaxations

- ▶ Easy to design
- ▶ Tractable biconjugation via **linear programming (LP)**
- ▶ Applicable to various submodular and **non-submodular** structures

Template for TU structures

Sparsity and structure together [5]

Given some weights $d \in \mathbb{R}^d$, $e \in \mathbb{R}^p$ and an integral vector $c \in \mathbb{Z}^l$, we define

$$g_{TU}(\mathbf{x}) := \min_{\omega} \{d^T \omega + e^T s : M \begin{bmatrix} \omega \\ s \end{bmatrix} \leq c, \mathbf{1}_{\text{supp}(\mathbf{x})} = s, \omega \in \{0, 1\}^d\}$$

for all feasible \mathbf{x} , ∞ otherwise. The parameter ω is useful for latent modeling.

Total unimodular (TU): $M \in \mathbb{R}^{l \times m}$ is TU iff the determinant of every square submatrix of M is 0, or ± 1 .

Relaxation of ILP to LP [10]

When M is TU and c is integral, then the LP

$$\max_{\beta \in \mathbb{R}^m} \{\theta^T \beta : M\beta \leq c, \beta \geq 0\}$$

has integer optimal solutions (i.e., ILP \equiv LP).

- ▶ “Exact convex relaxation” of: $g^*(\mathbf{y}) = \sup_{s \in \{0,1\}^p} |\mathbf{y}|^T s - F(s)$.
- ▶ Same idea behind the tractable biconjugation of submodular functions

Convexification of TU structures

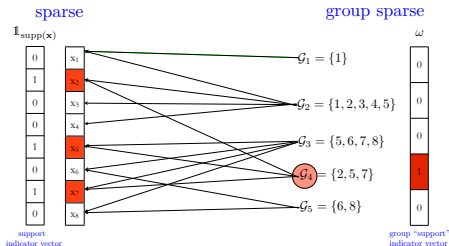
TU convex relaxation given by LP

$$g_{TU}^{**}(\mathbf{x}) := \min_{\omega} \{ \mathbf{d}^T \omega + \mathbf{e}^T \mathbf{s} : M \begin{bmatrix} \omega \\ \mathbf{s} \end{bmatrix} \leq \mathbf{c}, |\mathbf{x}| \leq \mathbf{s}, \omega \in \{0, 1\}^d \}$$

for all feasible \mathbf{x} , ∞ otherwise.

- ▶ Special cases:
 - ▶ Rederive the convex envelope of several submodular models
 - ▶ Establish the tightness of some convex regularizers for non-submodular models
- ▶ Beyond linear objectives, some quadratic objectives can also be handled

Group cover sparsity: Minimal group cover [2, 12, 8]



Structure: We seek the signal covered by a minimal number of groups.

$$\text{Objective: } d^T \omega$$

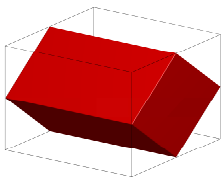
Linear description: For each non-zero coefficient, at least one group containing it is selected

$$B\omega \geq s$$

where B is the biadjacency matrix of \mathcal{G} , i.e., $B_{ij} = 1$ iff i -th coefficient is in \mathcal{G}_j .

When B is an interval matrix, or \mathcal{G} has a *loopless* group intersection graph it is TU.

Group cover sparsity: **Minimal group cover** [2, 12, 8]



$\mathcal{G} = \{\{1, 2\}, \{2, 3\}\}$, unit group weights $d = \mathbf{1}$.

Structure: *We seek the signal covered by a minimal number of groups.*

Objective: $d^T \omega$

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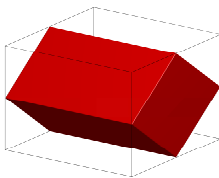
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Biconjugate: $g_{TU}^{**}(\mathbf{x}) = \min_{\omega \in [0,1]^M} \{d^T \omega : B\omega \geq |\mathbf{x}|\}$ for $\mathbf{x} \in [-1, 1]^p$, ∞ otherwise

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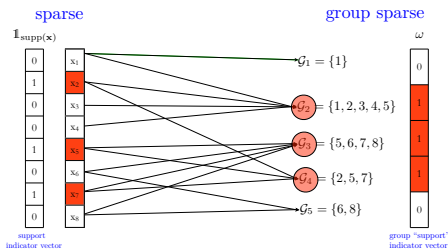
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 $\stackrel{*}{=} \min_{\mathbf{v}_i \in \mathbb{R}^p} \left\{ \sum_{i=1}^M d_i \|\mathbf{v}_i\|_{\infty} : \mathbf{x} = \sum_{i=1}^M \mathbf{v}_i, \forall \text{supp}(\mathbf{v}_i) \subseteq \mathcal{G}_i \right\},$

Group intersection sparsity [9, 14, 1]



Structure: We seek the signal intersecting with minimal number of groups.

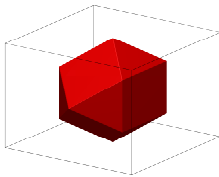
Objective: $d^T \omega$ (submodular: $F(\mathcal{S}) = \sum_{\mathcal{G}_i \in \mathcal{G}, \mathcal{S} \cap \mathcal{G}_i \neq \emptyset} d_i$)

Linear description: All groups containing a non-zero coefficient are selected

$$\mathbf{H}_k \mathbf{s} \leq \omega, \forall k \in \{0, \dots, p\}$$

where $\mathbf{H}_k(i, j) = \begin{cases} 1 & \text{if } j = k, j \in \mathcal{G}_i \\ 0 & \text{otherwise} \end{cases}$, which is TU.

Group intersection sparsity [9, 14, 1]



$$\mathfrak{G} = \{\{1, 2\}, \{2, 3\}\}, \text{ unit group weights } \mathbf{d} = \mathbf{1}$$

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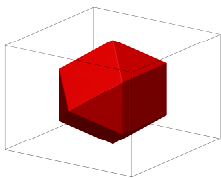
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for $\mathbf{x} \in [-1, 1]^p$, ∞ otherwise.

Group intersection sparsity [9, 14, 1]



$$\mathfrak{G} = \{\{1, 2\}, \{2, 3\}\}, \text{ unit group weights } \mathbf{d} = \mathbf{1}$$

Structure: *We seek the signal intersecting with minimal number of groups.*

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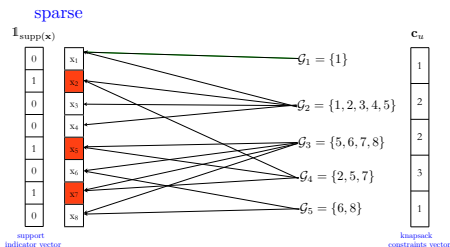
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for $\mathbf{x} \in [-1, 1]^p$, ∞ otherwise.

Group knapsack sparsity [15, 7, 6]



Structure: *We seek the sparsest signal with group allocation constraints.*

Objective: $\mathbb{1}^T \mathbf{s}$

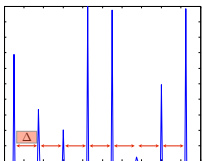
Linear description: A **valid** support obeys budget constraints over \mathbb{G}

$$\mathbf{B}^T \mathbf{s} \leq \mathbf{c}_u$$

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When \mathbf{B} is an interval matrix or \mathbb{G} has a *loopless* group intersection graph, it is **TU**.

Group knapsack sparsity [15, 7, 6]



$$B^T = \begin{bmatrix} 1 & 1 & \dots & 1 & 1 & 0 & 0 & \dots & 0 \\ 0 & 1 & 1 & \dots & 1 & 1 & 0 & \dots & 0 \\ & & & & & & & \ddots & \\ 0 & \dots & 0 & 0 & 1 & 1 & \dots & 1 & 1 \end{bmatrix}_{(p-\Delta+1) \times p}$$

Structure: *We seek the sparsest signal with group allocation constraints.*

Objective: $\mathbb{1}^T \mathbf{s}$

Linear description: A **valid** support obeys budget constraints over \mathcal{G}

$$B^T \mathbf{s} \leq \mathbf{c}_u$$

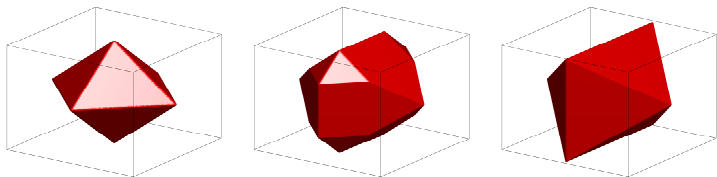
where B is the biadjacency matrix of \mathcal{G} , i.e., $B_{ij} = 1$ iff i -th coefficient is in \mathcal{G}_j .

When B is an interval matrix or \mathcal{G} has a **loopless** group intersection graph, it is **TU**.

Biconjugate: $g_{TU}^{**}(\mathbf{x}) = \begin{cases} \|\mathbf{x}\|_1 & \text{if } \mathbf{x} \in [-1, 1]^p, B^T |\mathbf{x}| \leq \mathbf{c}_u, \\ \infty & \text{otherwise} \end{cases}$

For the neuronal spike example, we have $\mathbf{c}_u = \mathbb{1}$.

Group knapsack sparsity [15, 7, 6]



(left) $g_{TU}^{**}(\mathbf{x}) \leq 1$ (middle) $g_{TU}^{**}(\mathbf{x}) \leq 1.5$ (right) $g_{TU}^{**}(\mathbf{x}) \leq 2$ for $\mathcal{G} = \{\{1, 2\}, \{2, 3\}\}$

Structure: *We seek the sparsest signal with group allocation constraints.*

Objective: $\mathbf{1}^T \mathbf{s}$

Linear description: A **valid** support obeys budget constraints over \mathcal{G}

$$\mathbf{B}^T \mathbf{s} \leq \mathbf{c}_u$$

where \mathbf{B} is the biadjacency matrix of \mathcal{G} , i.e., $B_{ij} = 1$ iff i -th coefficient is in \mathcal{G}_j .

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For the neuronal spike example, we have $\mathbf{c}_u = \mathbf{1}$.

Group knapsack sparsity example: A stylized spike train

- ▶ Basis pursuit (BP): $\|\mathbf{x}\|_1$
- ▶ TU-relax (TU):

$$g_{TU}^{**}(\mathbf{x}) = \begin{cases} \|\mathbf{x}\|_1 & \text{if } \mathbf{x} \in [-1, 1]^p, B^T|\mathbf{x}| \leq \mathbf{c}_u \\ \infty & \text{otherwise} \end{cases}$$

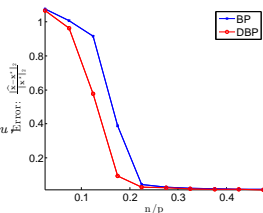
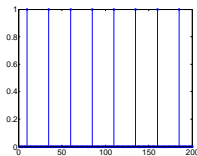
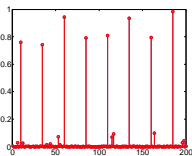


Figure: Recovery for $n = 0.18p$.



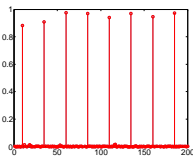
$\hat{\mathbf{x}}$

relative errors:



\mathbf{x}_{BP} solution

$$\frac{\|\hat{\mathbf{x}} - \mathbf{x}_{BP}\|_2}{\|\hat{\mathbf{x}}\|_2} = .200$$



\mathbf{x}_{TU} solution

$$\frac{\|\hat{\mathbf{x}} - \mathbf{x}_{TU}\|_2}{\|\hat{\mathbf{x}}\|_2} = .067$$

Conclusions

Our work: TU modeling framework

- ▶ Complement previous approaches
- ▶ Convex programs (not necessarily norms)
- ▶ Tight convexifications, non-submodular examples
- ▶ Easy to design and “usually” efficient via an LP

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